

TRUTHFUL MECHANISMS FOR GENERALIZED UTILITARIAN PROBLEMS

G. MELIDEO,¹ P.P.,² G. PROIETTI,^{1,3}
R. WATTENHOFER,⁴ P. WIDMAYER⁴

¹ UNIVERSITÀ DI L'AQUILA

² UNIVERSITÀ DI SALERNO

³ CNR

⁴ ETH ZÜRICH

INTERNET

- NO CENTRAL AUTHORITY
- SELF-INTERESTED COMPONENTS

- DIFFERENT GOALS

- NOT ALTRUISTIC

- MAY NOT FOLLOW THE "PROTOCOL"

PRIVATE COMPANIES

AUTONOMOUS SYSTEMS

PROVIDERS

UNIVERSITIES

⋮

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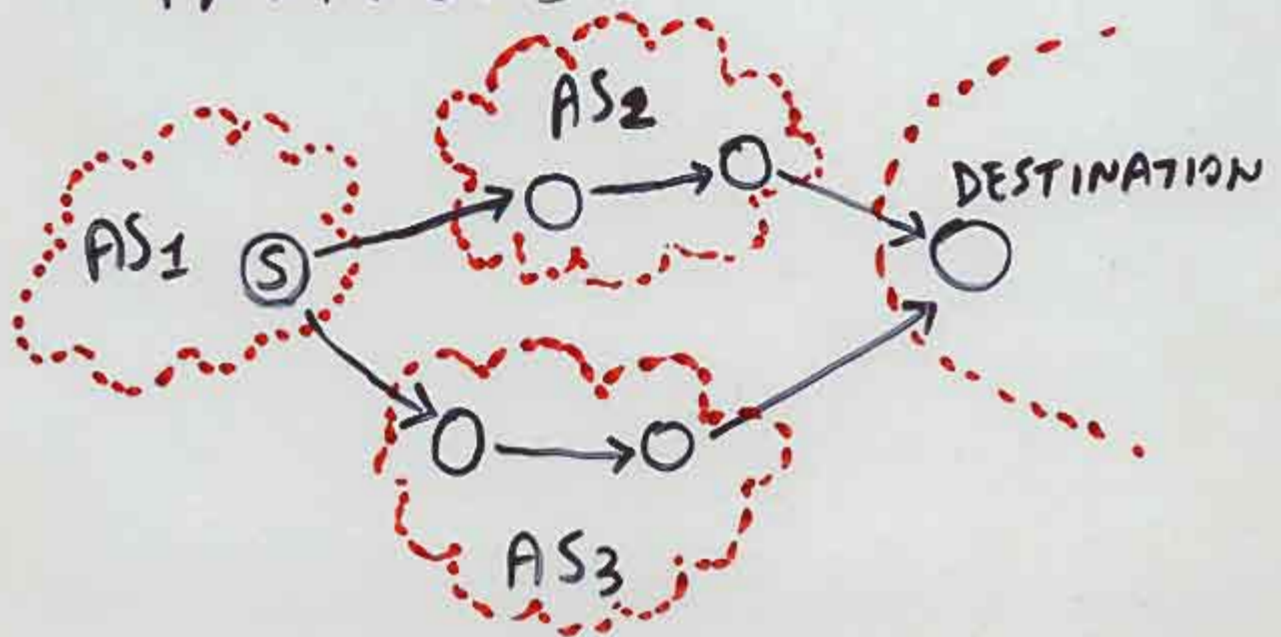


SELFISH AGENTS

INTERNET ROUTING

AUTONOMOUS SYSTEMS

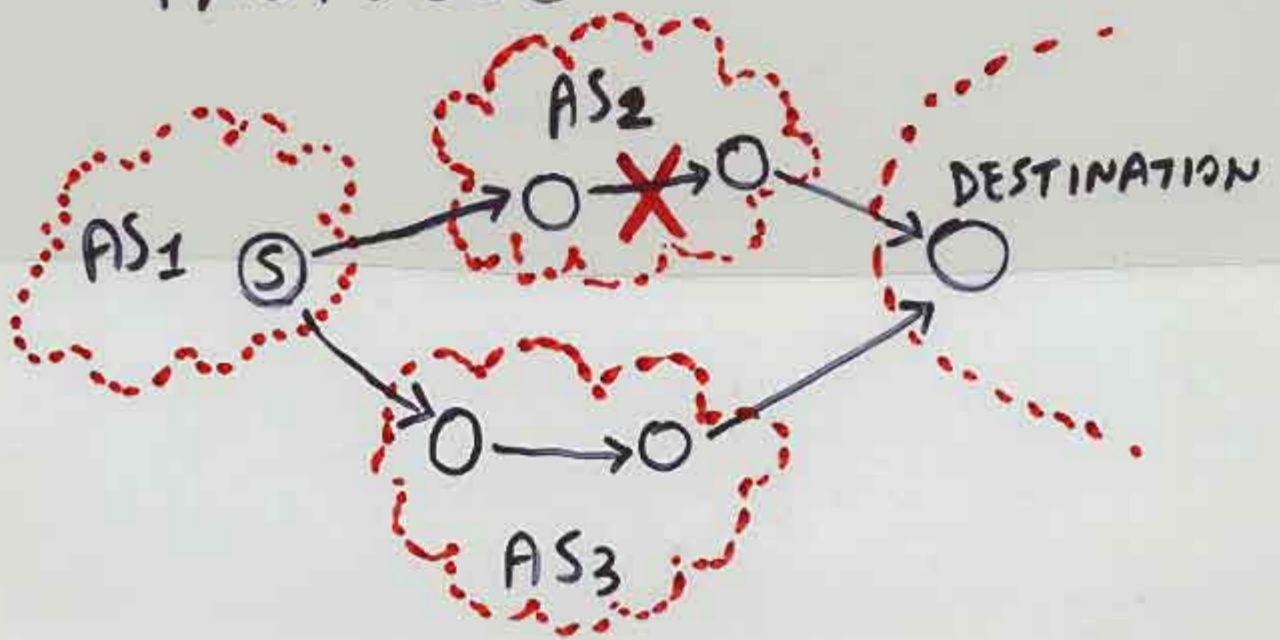
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- EXCHANGE INFORMATION
- IMPLEMENT ROUTING PROTOCOL



INTERNET ROUTING

AUTONOMOUS SYSTEMS

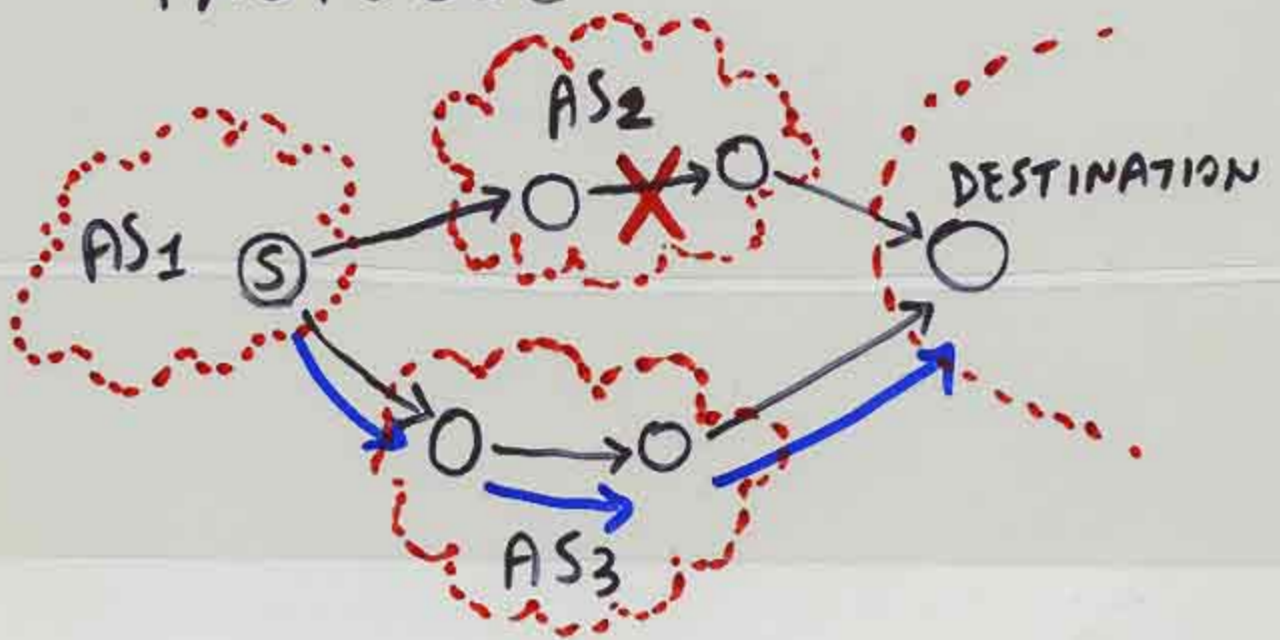
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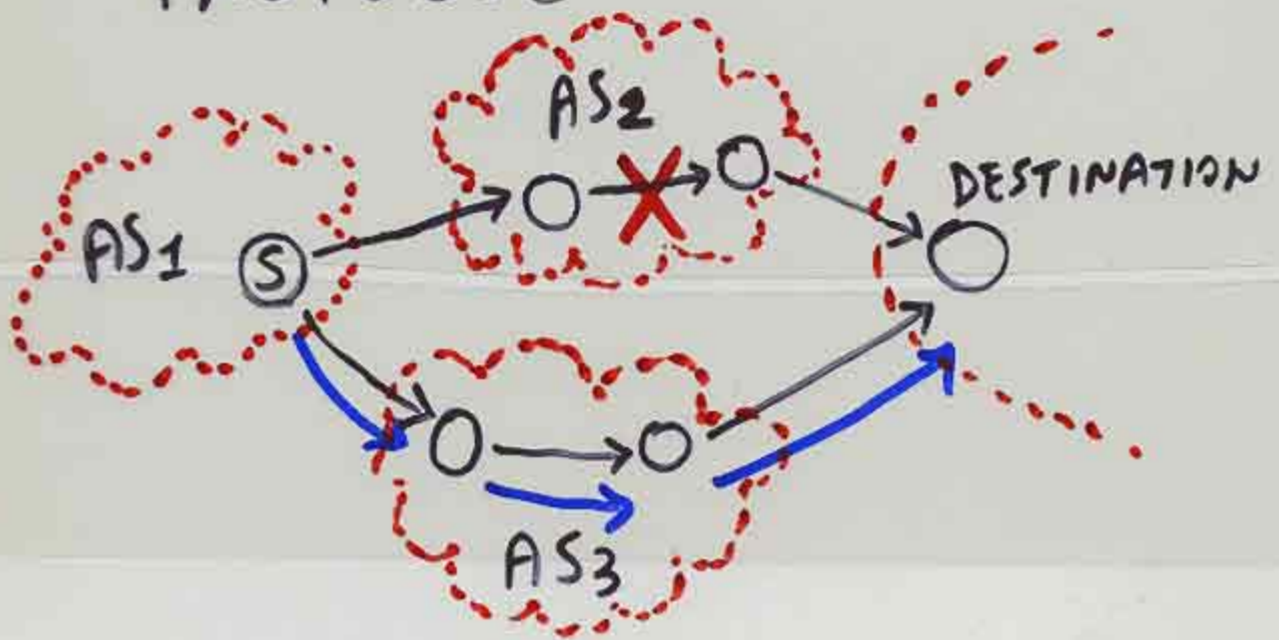
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INTERNET ROUTING

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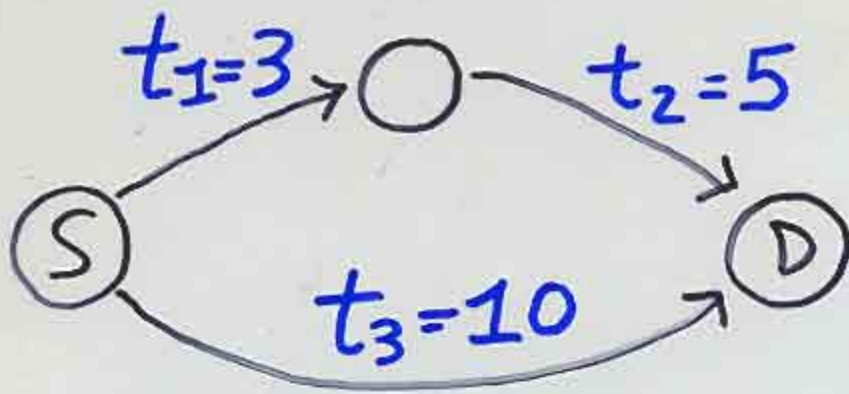
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FALSE LINK STATUS

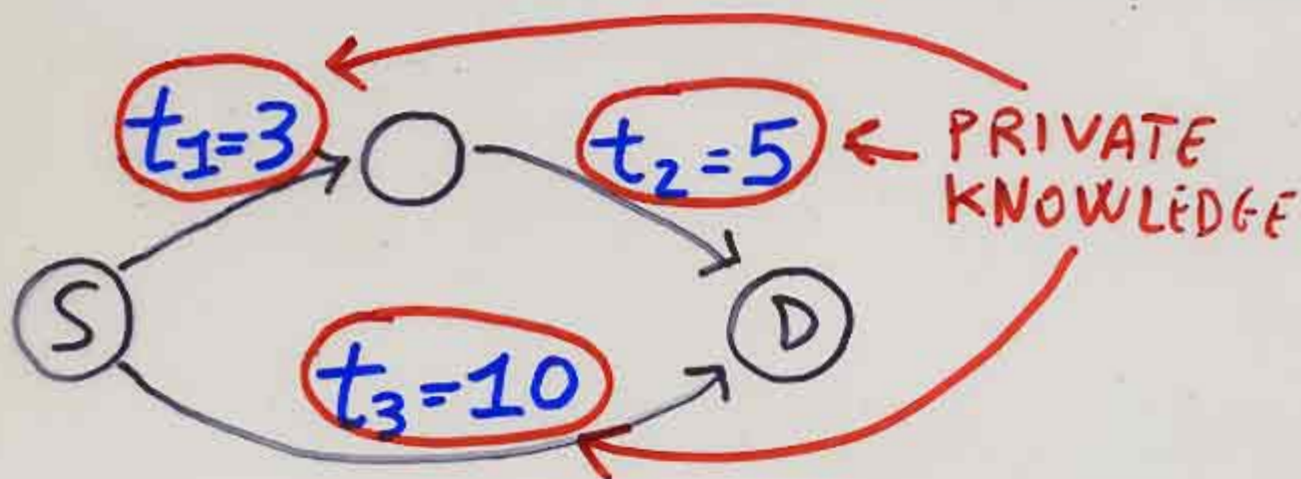


REDIRECT THE TRAFFIC



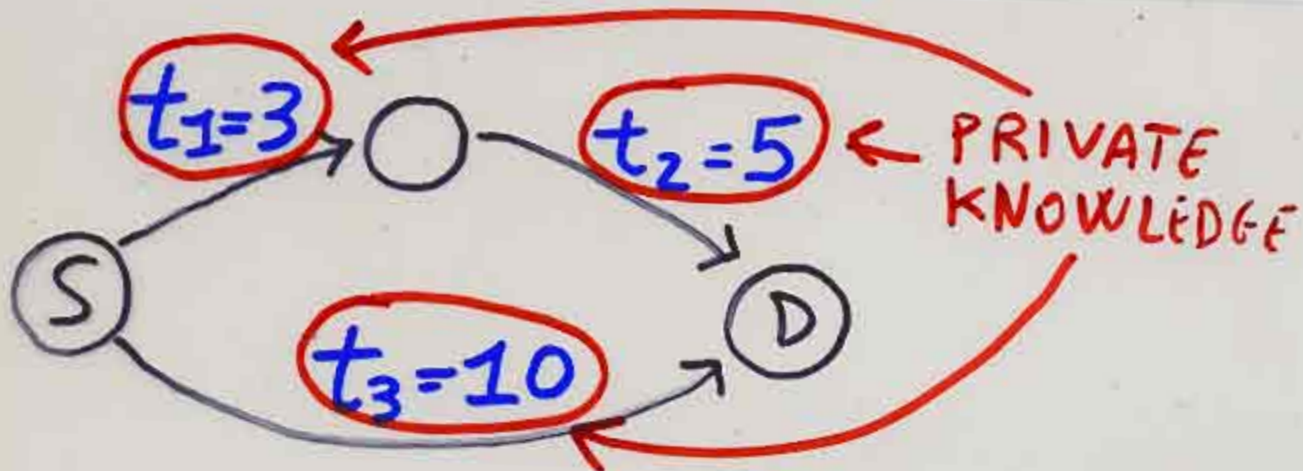
t_i = TIME TO PROCESS
1 PACKET OVER
LINK i

GOAL: PICK THE CHEAPEST
PATH FROM S TO D



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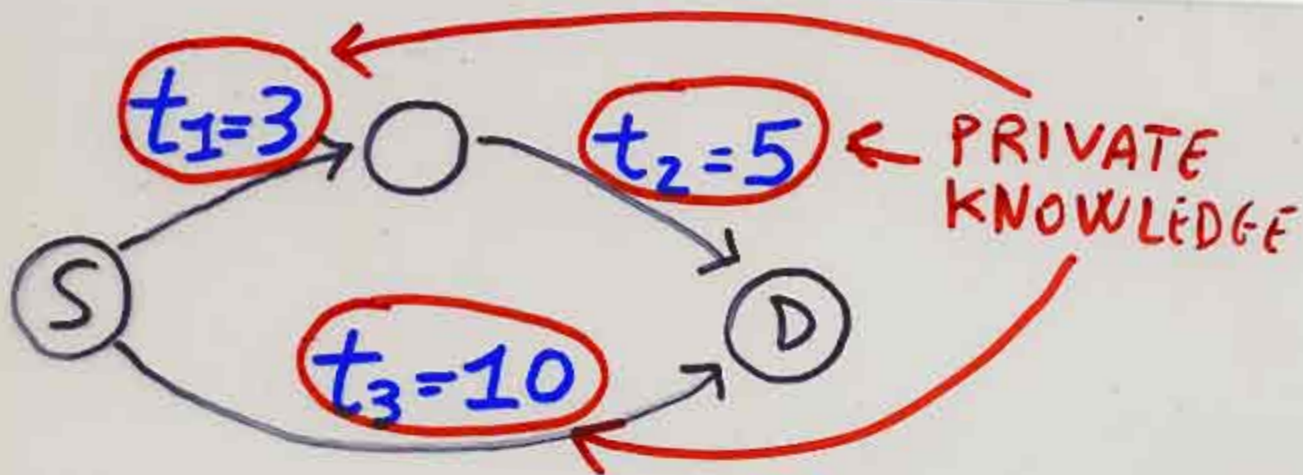


t_i = TIME TO PROCESS
1 PACKET OVER
LINK i



COST FOR AGENT i

GOAL: PICK THE CHEAPEST
PATH FROM S TO D



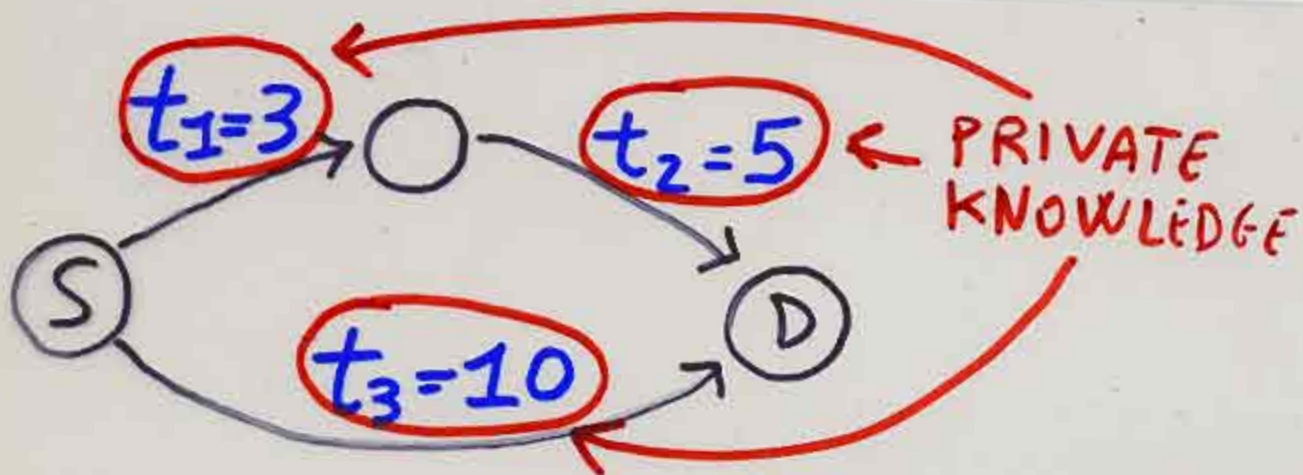
t_i = TIME TO PROCESS
1 PACKET OVER
LINK i



COST FOR AGENT i

GOAL: PICK THE CHEAPEST
PATH FROM S TO D

AGENTS CAN CHEAT:
REPORT $\pi_i \neq t_i$



t_i = TIME TO PROCESS
1 PACKET OVER
LINK i



COST FOR AGENT i

GOAL: PICK THE CHEAPEST
PATH FROM S TO D

AGENTS CAN CHEAT:
REPORT $\tau_i \neq t_i$

GIVE INCENTIVES!

PROBLEMS INVOLVING SELFISH AGENTS

INPUT: $I = (\pi, t)$

π → PUBLIC

t → PRIVATE

t_1, \dots, t_m

AG_1 AG_m

MEASURE: SOLUTION X

$$m(X, t)$$

DEPENDS ON PRIVATE INPUT

GOAL: OPTIMIZE $m(X, t)$
(MAXIMIZE)

PROBLEMS INVOLVING SELFISH AGENTS

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AG_1 AG_m

MEASURE: SOLUTION X

$$m(X, t)$$

DEPENDS ON PRIVATE INPUT

GOAL: OPTIMIZE $m(X, t)$
(MAXIMIZE)

WE NEED t_1, \dots, t_m !

MECHANISM DESIGN

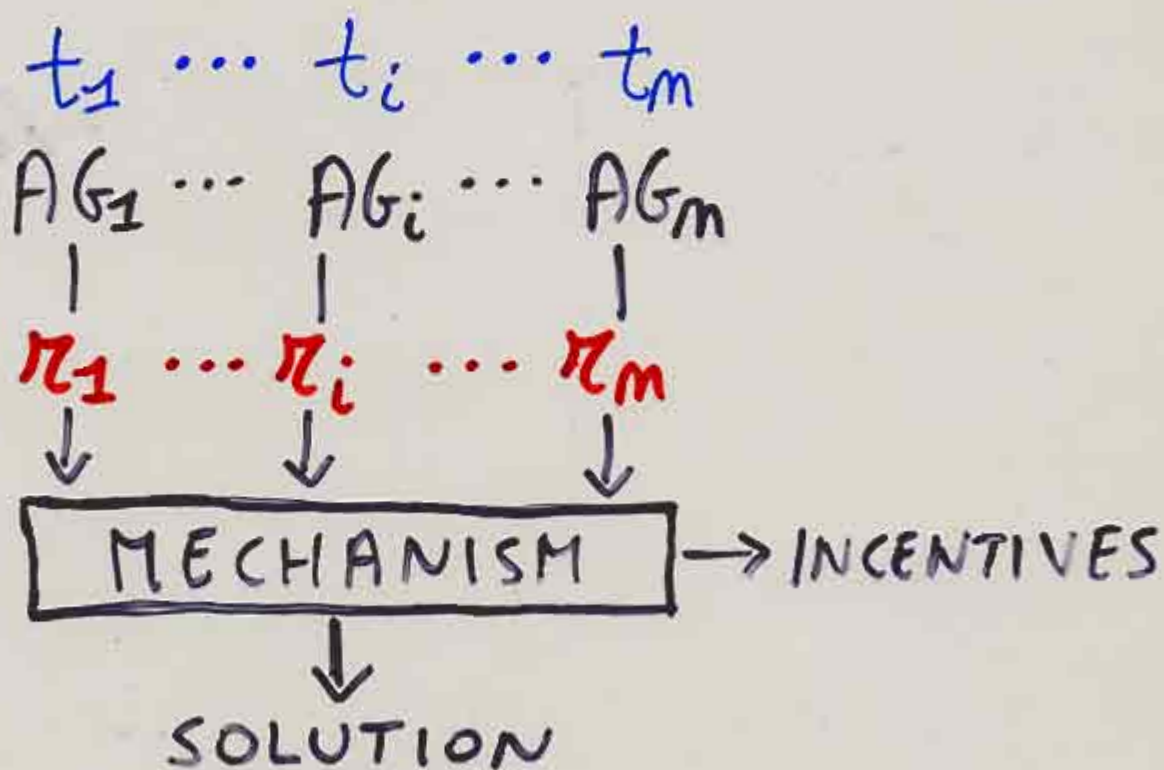
MECHANISM:

ALGORITHM + INCENTIVES
(e.g. PAYMENTS)

MECHANISM DESIGN

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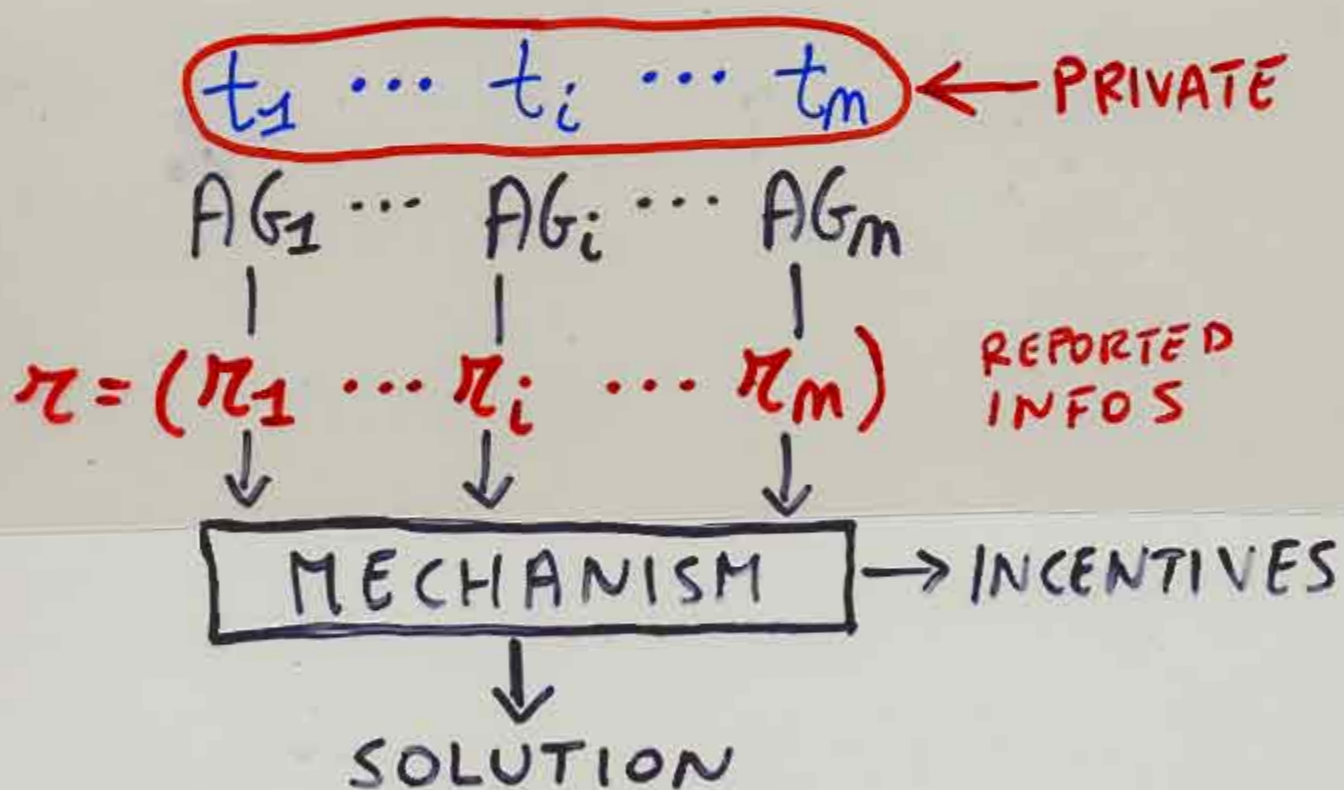
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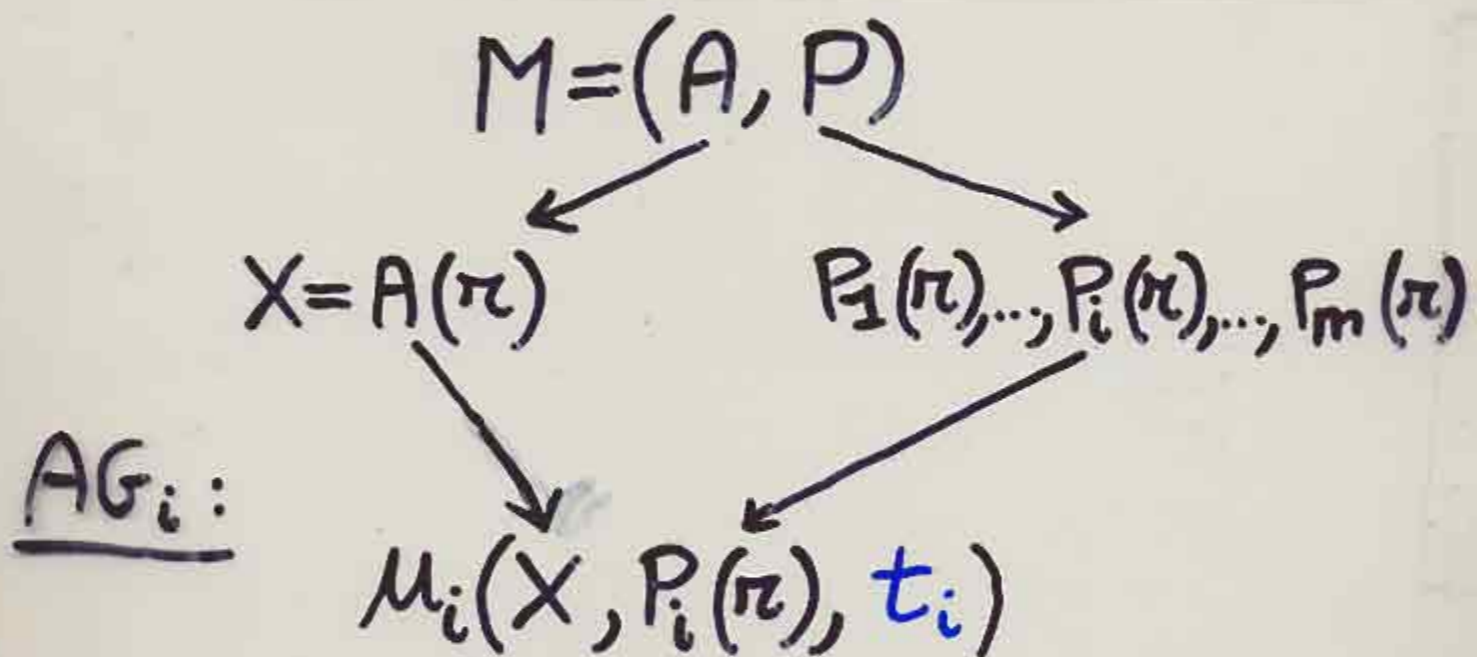
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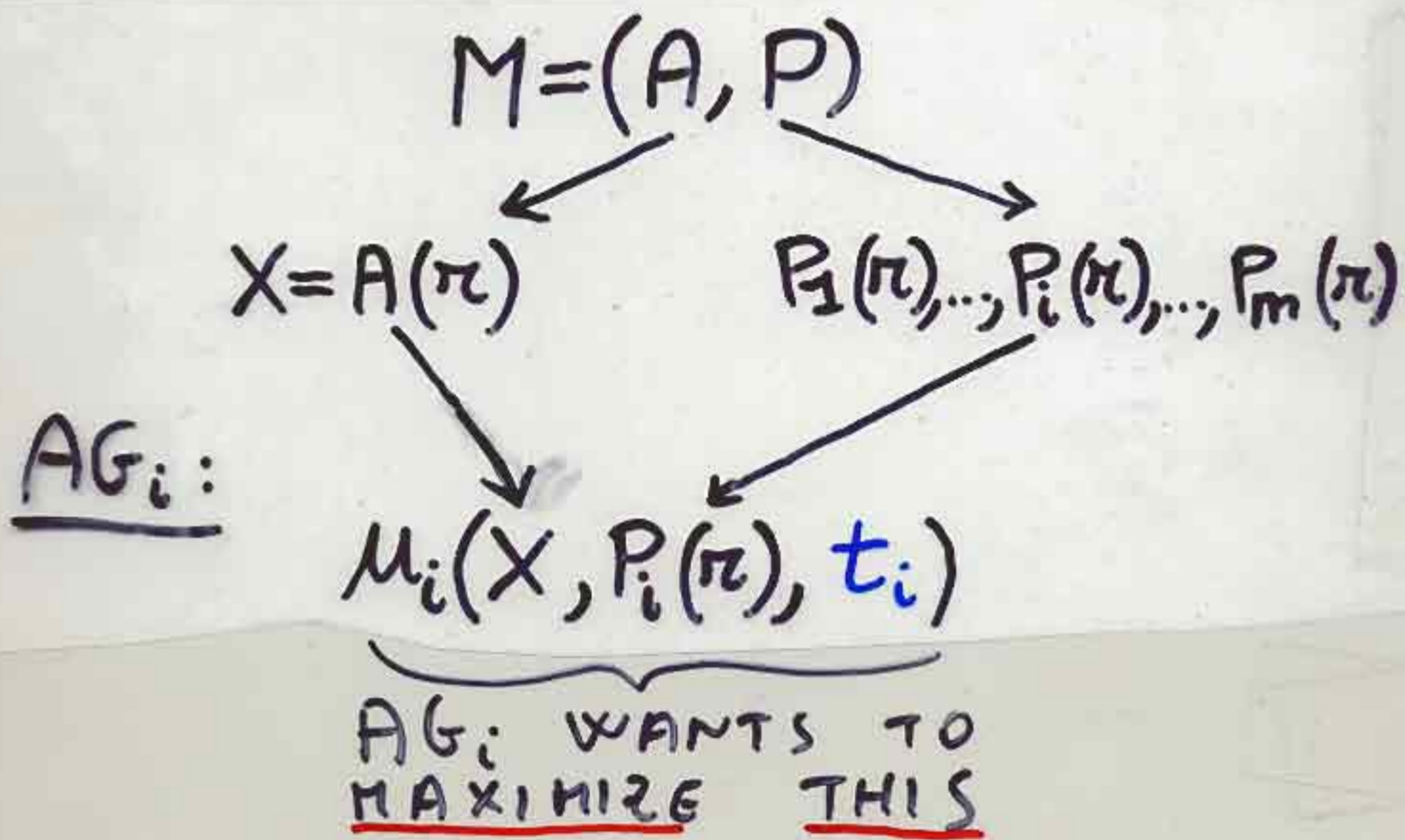
X AND $P_i(\pi)$ DEPEND ON

$\pi = (\pi_1, \dots, \pi_i, \dots, \pi_m)$

MECHANISM DESIGN

MECHANISM:

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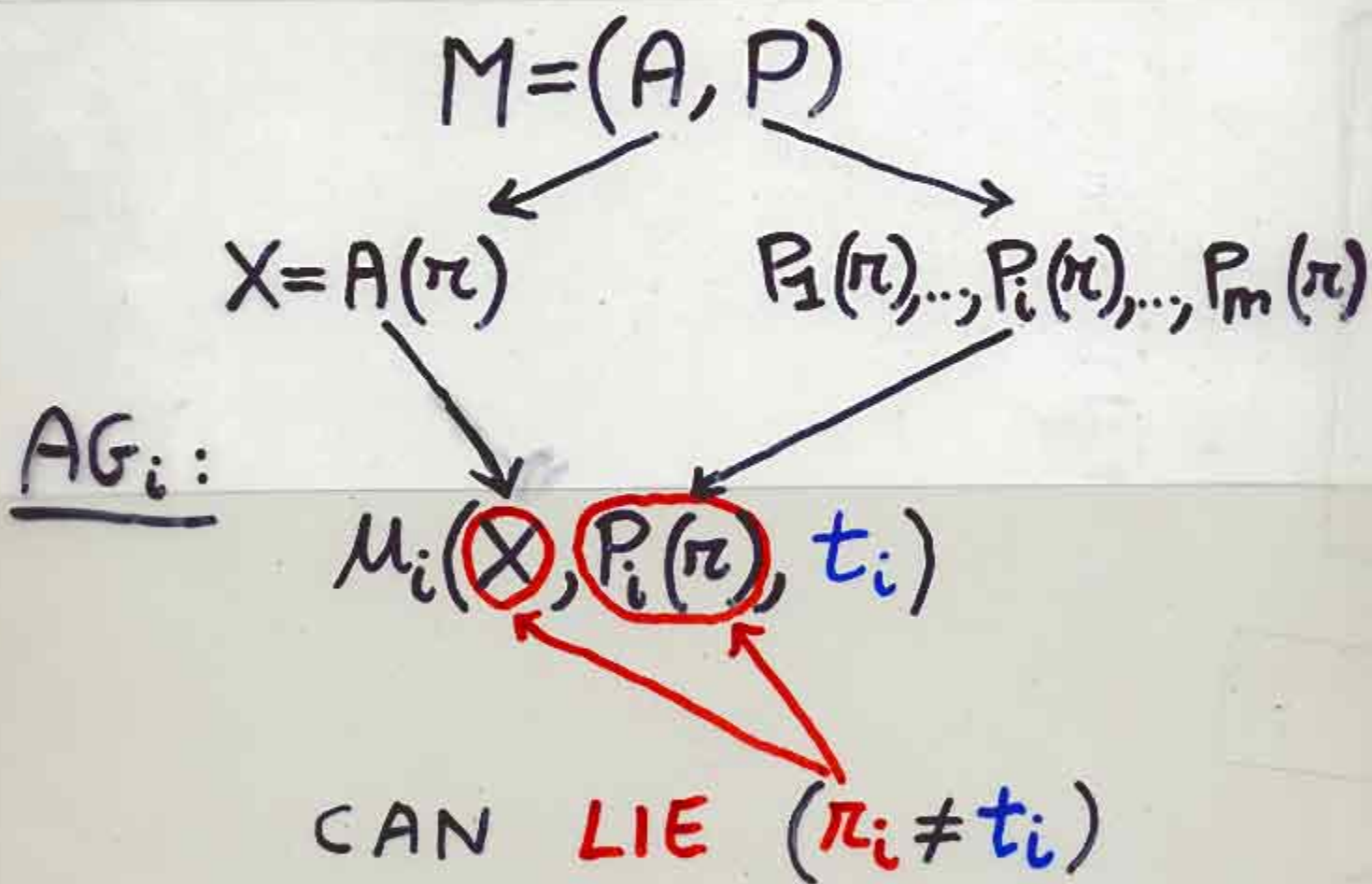


X AND $P_i(\pi)$ DEPEND ON
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MECHANISM DESIGN

MECHANISM:

ALGORITHM + INCENTIVES
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X AND $P_i(\pi)$ DEPEND ON

$$\pi = (\pi_1, \dots, \underline{\pi_i}, \dots, \pi_m)$$

MAIN ISSUES

WHICH PROBLEMS CAN
BE SOLVED?

- $M=(A,P)$ IS TRUTHFUL
- $A(t)$ IS OPTIMAL SOLUTION

GENERAL TECHNIQUES?

ANSWER:

UTILITARIAN PROBLEMS

CAN BE SOLVED WITH

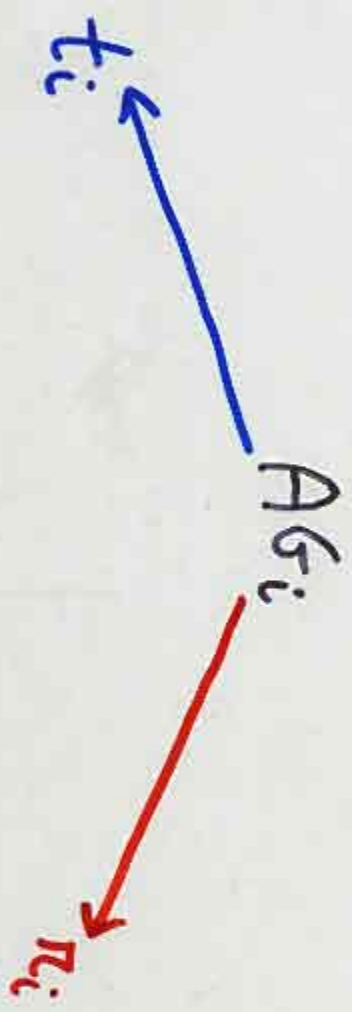
VCG MECHANISMS

[VICREY '61, CLARKE '71, GROVES '73]

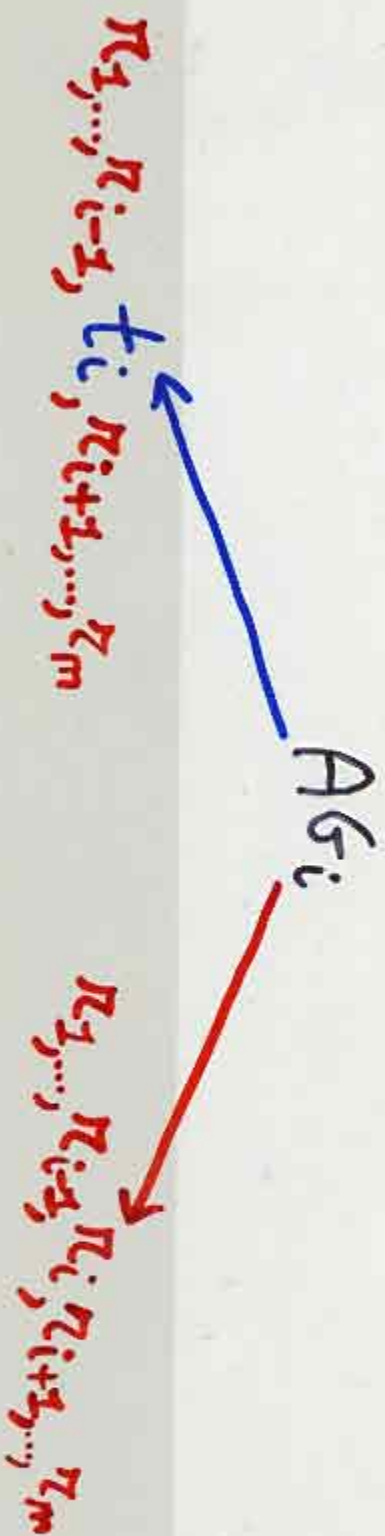
THE ONLY WAY (IN SOME CASES)

[GREEN-LAFFONT '77]

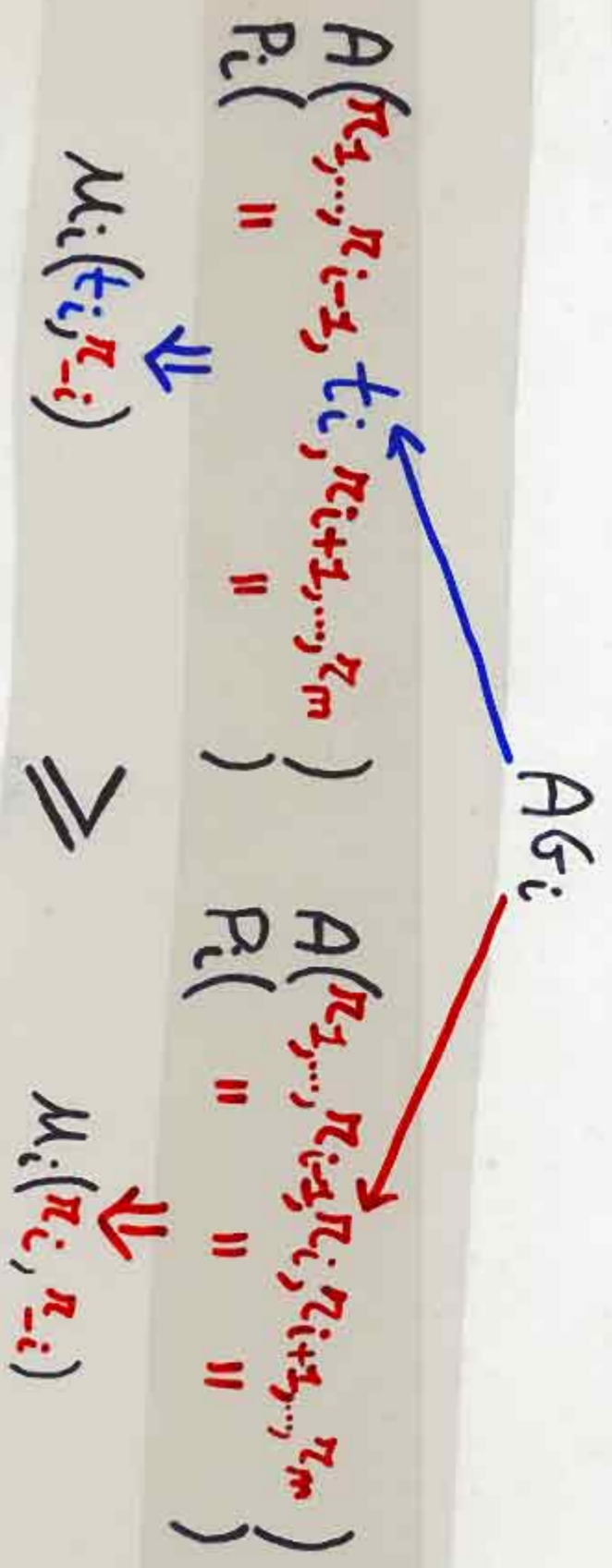
MECHANISM $M = (A, P)$ IS TRUTHFUL IF



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$$\forall i, \forall \pi_{-i} = (\pi_1, \dots, \pi_{i-1}, \pi_{i+1}, \dots, \pi_m), \forall \pi_i$$

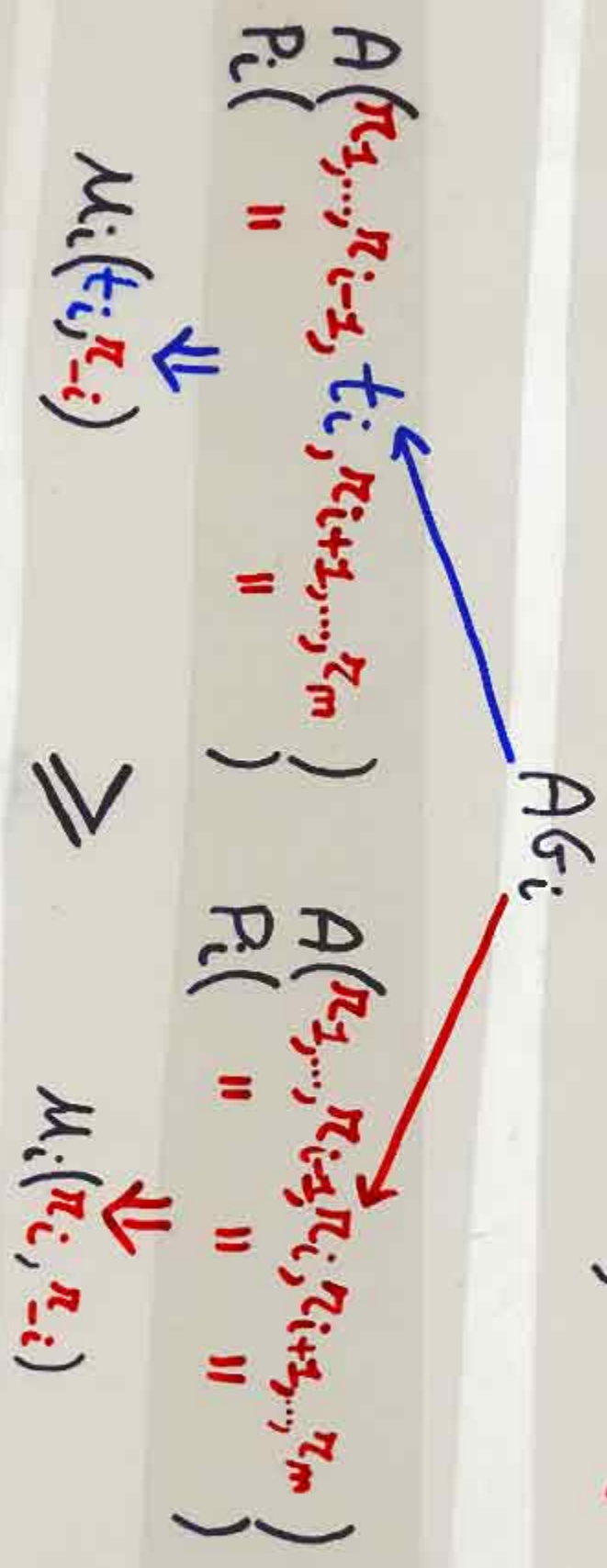
$A \subseteq_i$

$$A(\pi_1, \dots, \pi_{i-1}, t_i, \pi_{i+1}, \dots, \pi_m) \succeq_i A(\pi_1, \dots, \pi_{i-1}, \pi_i, \pi_{i+1}, \dots, \pi_m)$$

$$M_i(t_i, \pi_{-i}) \succeq_i M_i(\pi_i, \pi_{-i})$$

MECHANISM $M = (A, P)$ IS TRUTHFUL IFF

$$\forall i, \forall \pi_{-i} = (\pi_1, \dots, \pi_{i-1}, \pi_{i+1}, \dots, \pi_m), \forall \pi_i$$



NONE HAS AN INCENTIVE TO LIE, EVER !!

UTILITARIAN PROBLEMS

UTILITIES:

$$U_i(P_i, X, t_i)$$

||

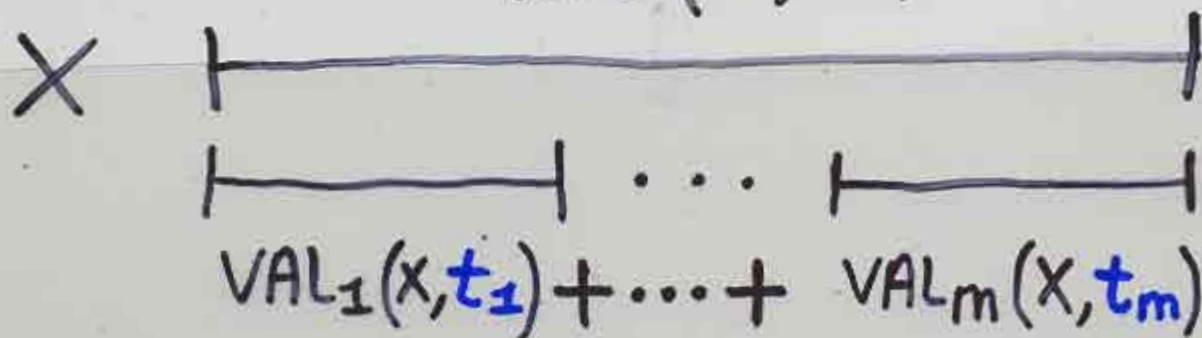
$$VAL_i(X, t_i) + P_i$$

↑
HOW MUCH
AG: "LIKES"
SOLUTION X

↑
MONEY
RECEIVED

MEASURE:

$$m(X, t)$$



SOLUTIONS: $SOL(I) = SOL(\pi)$

↑
PUBLIC

UTILITARIAN PROBLEMS

UTILITIES:

$$U_i(P_i, X, t_i)$$

||

$$VAL_i(X, t_i) + P_i$$

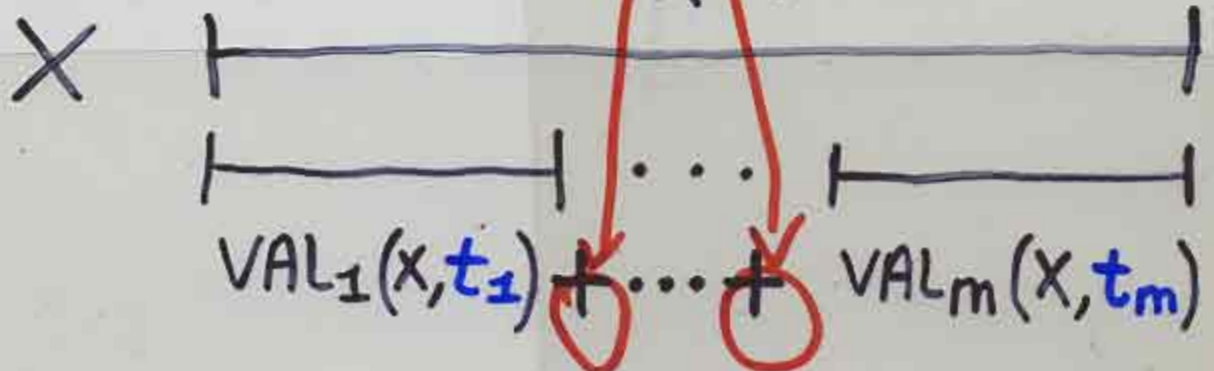
TOO RESTRICTIVE

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HOW MUCH
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MEASURE:

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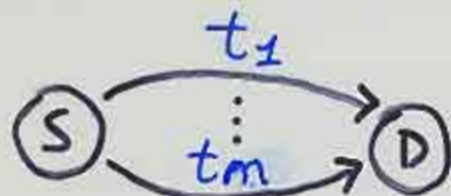
TOO RESTRICTIVE

SOLUTIONS: $SOL(I) = SOL(\pi)$

↑
PUBLIC

NON-UTILITARIAN CASES

EX1: MOST RELIABLE LINK



$$t_i = P_R [\text{LINK } i \text{ DOES NOT FAIL}]$$

INCENTIVE: PAY P_i IF
LINK i SUCCEEDED

$$\text{UTILITY: } U_i(X, P_i, t_i) = P_i \cdot t_i$$

\downarrow
 $VAL_i(X, t_i)$

EX2: TASK SCHEDULING

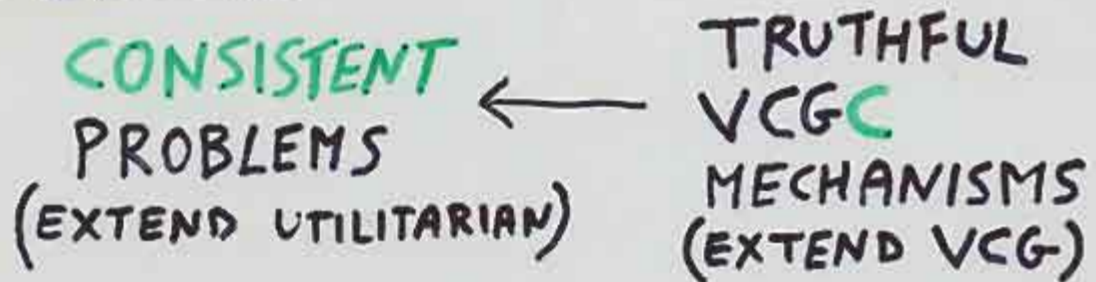
X = JOB ALLOCATION

$VAL_i(X, t_i)$ = FINISH TIME OF
MACHINE i

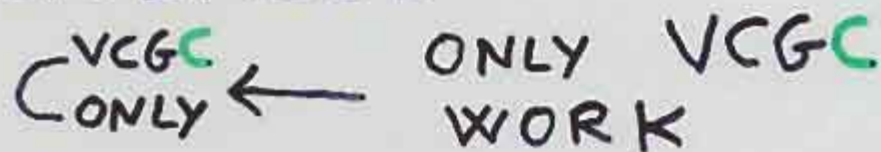
$$m(X, t) = \text{MAX}(VAL_1(X, t_1), \dots, VAL_m(X, t_m))$$

OUR CONTRIBUTION

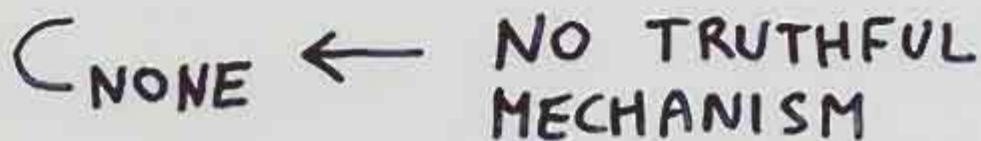
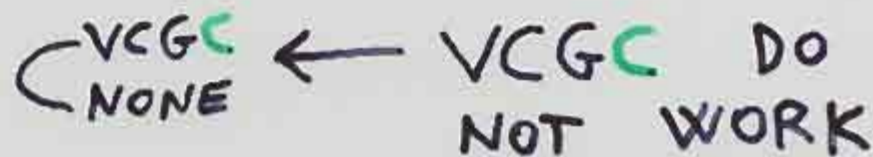
① EXTEND THE VCG IDEA TO A WIDER CLASS OF PROBLEMS:



② UNIQUENESS



③ LIMITS



④ APPLICATIONS

NON-UTILITARIAN PROBLEMS

(BASIC CS PROBLEMS)

INTUITION

OUR GOAL

$m(x, t)$

AG_i's GOAL

$\mu_i(x, P_i, t_i)$

POSSIBLE "=" IF $m(\cdot)$ AND $\mu_i(\cdot)$
HAVE A COMMON "ROOT"



INTUITION

OUR GOAL

AG_i's GOAL

$$m(x, t) = M_i(x, p_i, t_i)$$

AG_i IS WILLING TO HELP US
(MAXIMIZE $m(x, t)$)



1. We have a common goal



INTUITION

OUR GOAL

AG_i 's GOAL

$$m(x, t) = M_i(x, P_i, t_i)$$

AG_i IS WILLING TO HELP US

(MAXIMIZE $m(x, t)$)



POSSIBLE "=" IF $m(\cdot)$ AND $M_i(\cdot)$
HAVE A COMMON "ROOT"

INTUITION

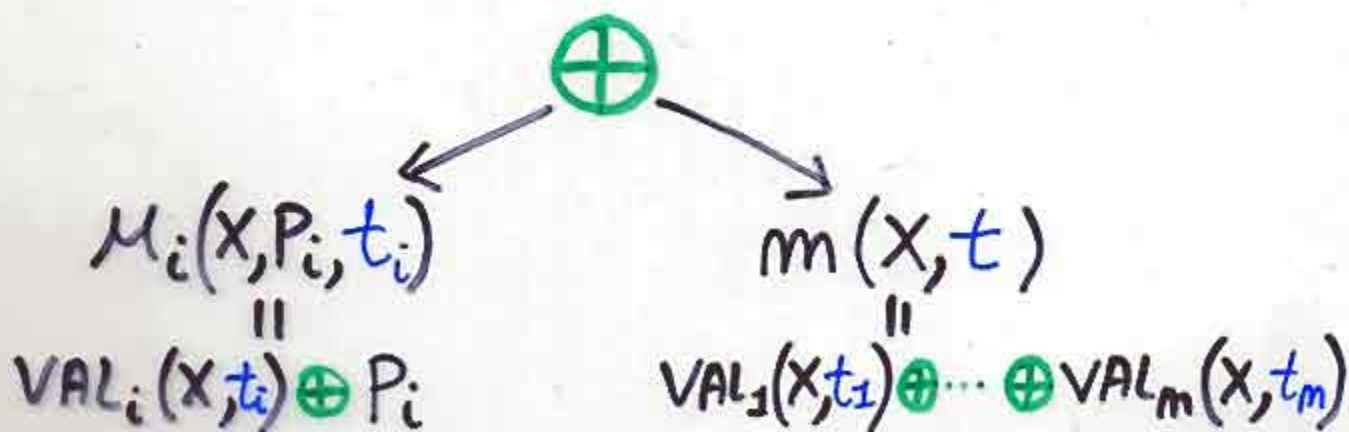
OUR GOAL

AG_i's GOAL

$$m(x, t) = \mu_i(x, P_i, t_i)$$

AG_i IS WILLING TO HELP US
(MAXIMIZE $m(x, t)$) 😊

POSSIBLE "=" IF $m(\cdot)$ AND $\mu_i(\cdot)$
HAVE A COMMON "ROOT"



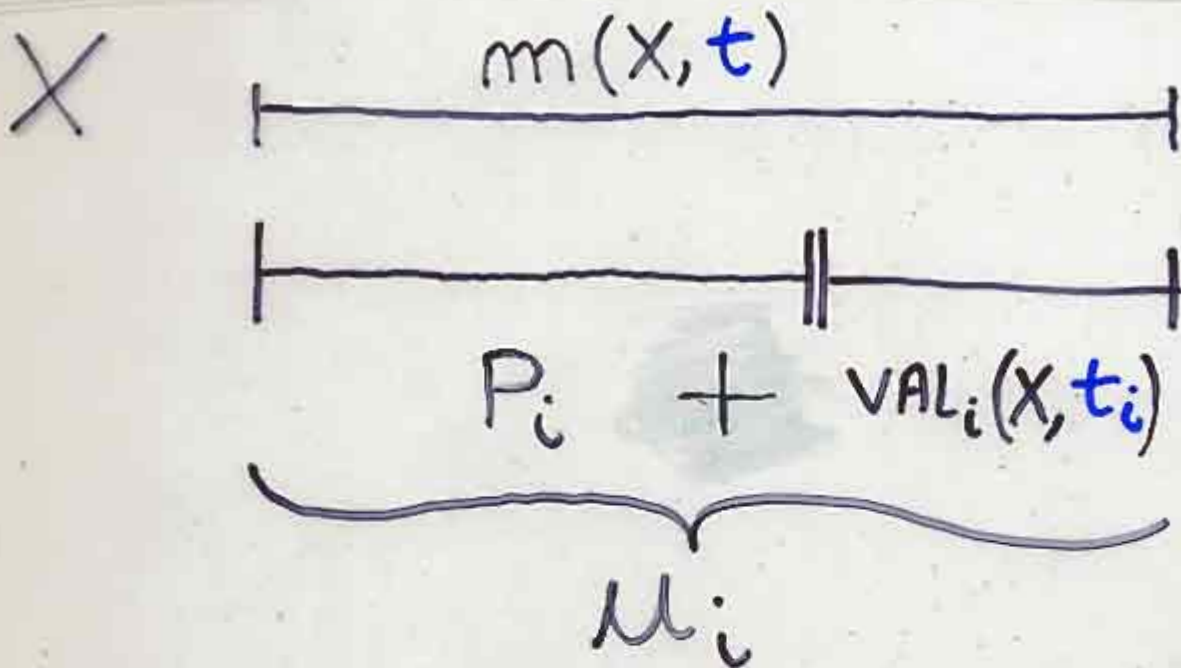
VCG PAYMENTS

KNOW ALL BUT ONE (t_i)



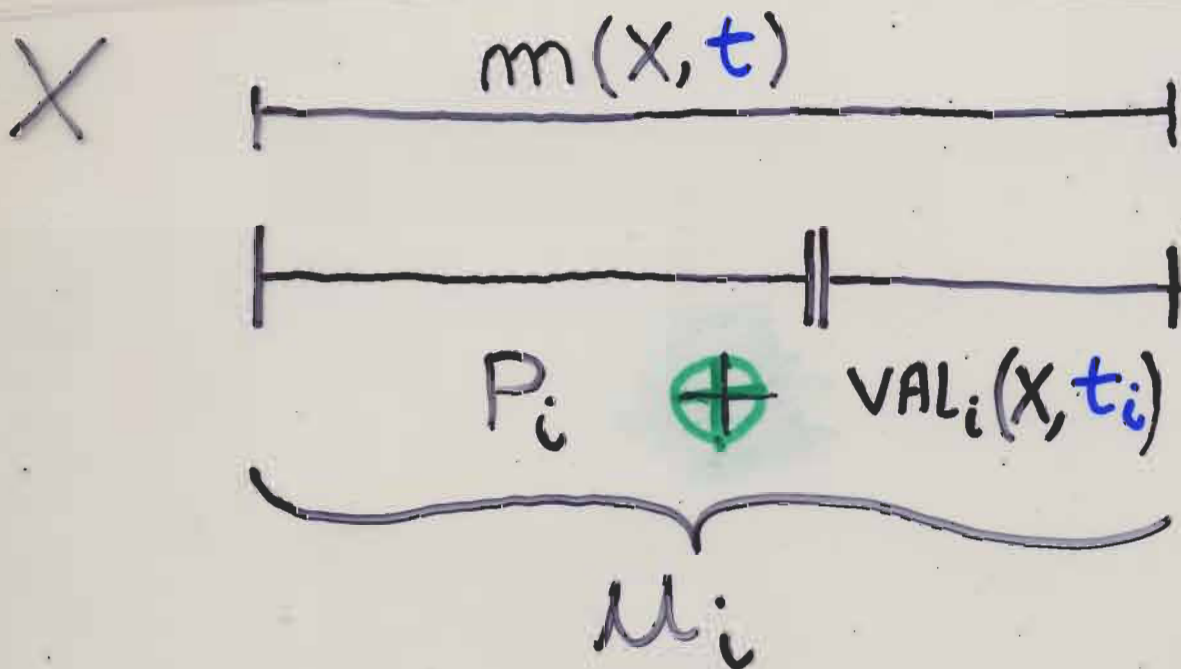
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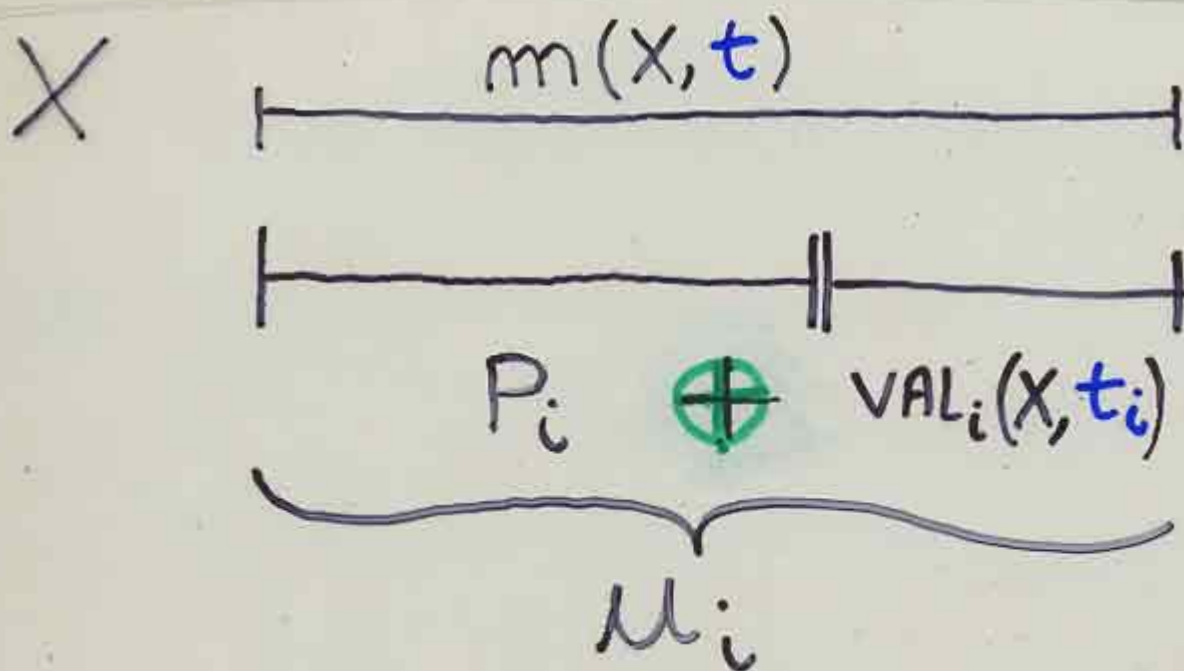
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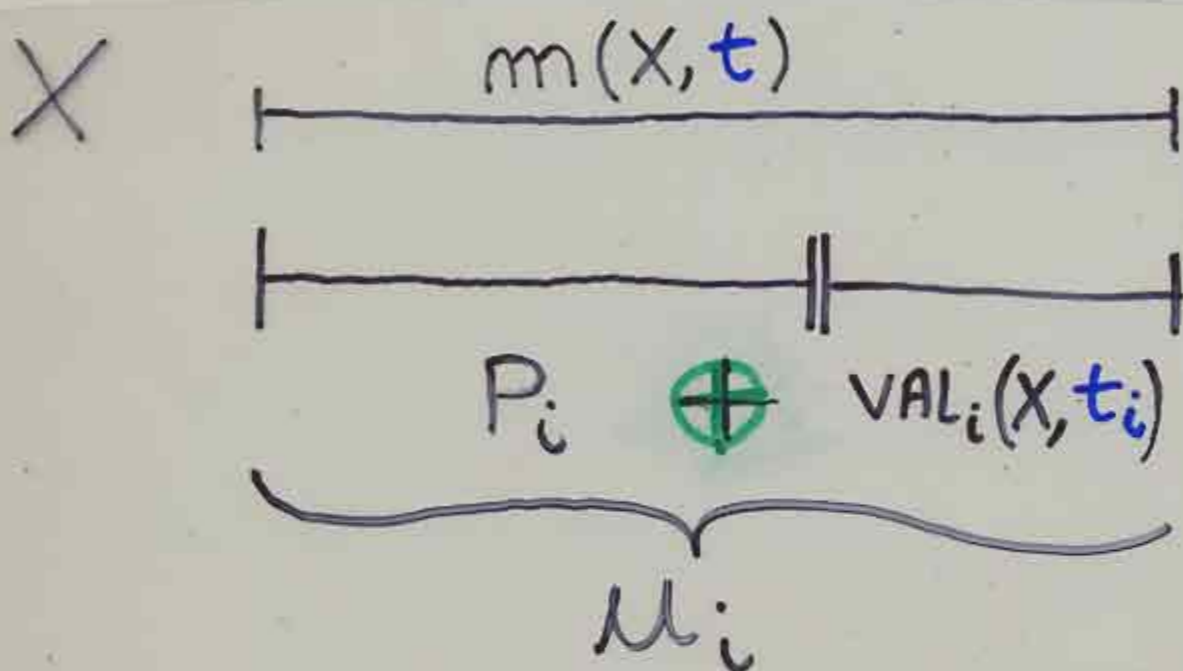


$$P_i = \bigoplus_{j \neq i} \text{VAL}_j(X, \tau_j)$$

\Downarrow

$$\mu_i = m(X, (\tau_1, \dots, \tau_{i-1}, t_i, \tau_{i+1}, \dots, \tau_m))$$

VCG PAYMENTS

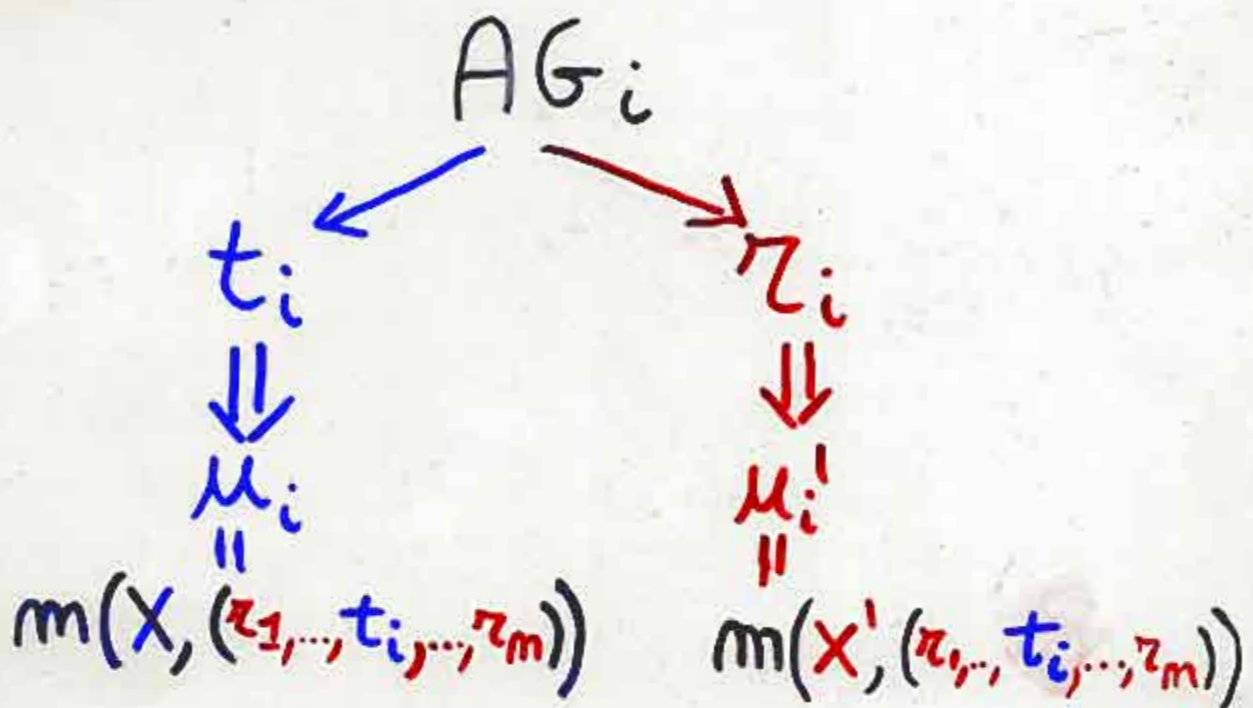


$$P_i = \bigoplus_{j \neq i} \text{VAL}_j(X, \pi_j)$$

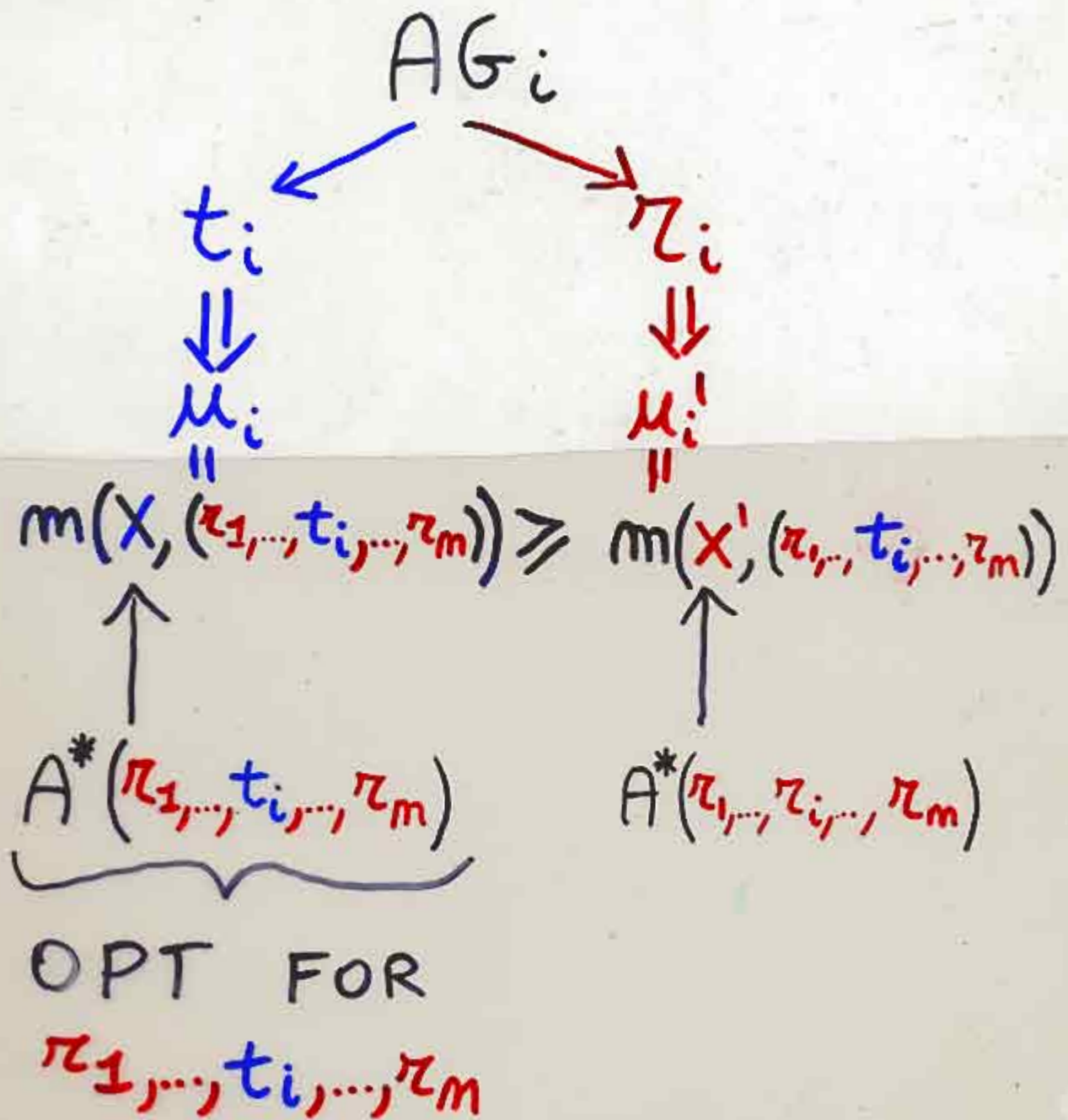
\Downarrow ? ('+' OK)

$$u_i = m(X, (\pi_1, \dots, \pi_{i-1}, t_i, \pi_{i+1}, \dots, \pi_m))$$

TRUTHFULNESS:



TRUTHFULNESS:



$$\mu_i(X, P_i, t_i) \stackrel{?}{\Rightarrow} m(X, \cdot)$$

DICOMPOSE $m(X, \pi)$ [PAYMENTS]

$$VAL_1(X, \pi_1) \oplus \dots \oplus VAL_i(X, \pi_i) \oplus \dots \oplus VAL_m(X, \pi_m)$$

RIASSEMBLE $m(X, \cdot)$ [UTILITY]

$$\begin{aligned} \mu_i(X, P_i, t_i) &= P_i \oplus VAL_i(X, t_i) = \\ & (A \oplus B) \oplus VAL_i(X, t_i) \end{aligned}$$

$$\mu_i(X, P_i, t_i) \stackrel{?}{\Rightarrow} m(X, \cdot)$$

DECOMPOSE $m(X, \pi)$ [PAYMENTS]

$$\text{VAL}_1(X, \pi_1) \oplus \dots \oplus \text{VAL}_i(X, \pi_i) \oplus \dots \oplus \text{VAL}_m(X, \pi_m)$$

DECOMPOSE $m(X, \cdot)$ [PAYMENTS]

$$\mu_i(X, P_i, t_i) = P_i \oplus \text{VAL}_i(X, t_i)$$

$$(A \oplus B) \oplus \text{VAL}(X, t)$$

$$\mu_i(X, P_i, t_i) \stackrel{?}{\Rightarrow} m(X, \cdot)$$

DI-COMPOSE

$m(X, \tau)$ [PAYMENTS]

$$\underbrace{\text{VAL}_1(X, \tau_1) \oplus \dots \oplus \text{VAL}_i(X, \tau_i) \oplus \dots \oplus \text{VAL}_m(X, \tau_m)}_{A \oplus B}$$

A

B

$$P_i = A \oplus B$$

EXAMPLE $m(X, \tau)$ [PAYMENTS]

$$\mu_i(X, P_i, t_i) = P_i \oplus \text{VAL}_i(X, t_i) = \dots$$

$$(A \oplus B) \oplus \text{VAL}_i(X, t_i)$$

$$\mu_i(X, P_i, t_i) \stackrel{?}{\Rightarrow} m(X, \cdot)$$

DICOMPOSE $m(X, \pi)$ [PAYMENTS]

$$\underbrace{VAL_1(X, \pi_1) \oplus \dots \oplus VAL_i(X, \pi_i) \oplus \dots \oplus VAL_m(X, \pi_m)}_{A \oplus B}$$

$P_i = A \oplus B$

RIASSEMBLE $m(X, \cdot)$ [UTILITY]

$$\mu_i(X, P_i, t_i) = P_i \oplus VAL_i(X, t_i) = (A \oplus B) \oplus VAL_i(X, t_i)$$

$$u_i(x, p_i, t_i) \stackrel{?}{\Rightarrow} m(x, \cdot)$$

DISCOMPOSE $m(x, \pi)$ [PAYMENTS]

$$\underbrace{VAL_1(x, \pi_1) \oplus \dots \oplus VAL_i(x, \pi_i)}_A \oplus \dots \oplus \underbrace{VAL_m(x, \pi_m)}_B$$

$P_i = A \oplus B$

RIASSEMBLE $m(x, \cdot)$ [UTILITY]

$$u_i(x, p_i, t_i) = P_i \oplus VAL_i(x, t_i) =$$

$$(A \oplus B) \oplus VAL_i(x, t_i)$$

||

$$A \oplus VAL_i(x, t_i) \oplus B$$

$$\mu_i(X, P_i, t_i) \stackrel{?}{\Rightarrow} m(X, \cdot)$$

DICOMPOSE $m(X, \pi)$ [PAYMENTS]

$$\underbrace{\text{VAL}_1(X, \pi_1) \oplus \dots \oplus \text{VAL}_i(X, \pi_i) \oplus \dots \oplus \text{VAL}_m(X, \pi_m)}_{A \oplus B}$$

$P_i = A \oplus B$

RIASSEMBLE $m(X, \cdot)$ [UTILITY]

$$\mu_i(X, P_i, t_i) = P_i \oplus \text{VAL}_i(X, t_i) =$$

$$(A \oplus B) \oplus \text{VAL}_i(X, t_i)$$

||

$$A \oplus \text{VAL}_i(X, t_i) \oplus B$$

$$m(X, (\pi_1, \dots, \pi_{i-1}, t_i, \pi_{i+1}, \dots, \pi_m))$$

$$\mu_i(X, P_i, t_i) \stackrel{?}{\Rightarrow} m(X, \cdot)$$

DI-COMPOSE $m(X, \pi)$ [PAYMENTS]

$$\underbrace{\text{VAL}_1(X, \pi_1) \oplus \dots \oplus \text{VAL}_i(X, \pi_i) \oplus \dots \oplus \text{VAL}_m(X, \pi_m)}_{A \oplus B}$$

$P_i = A \oplus B$

RI-ASSEMBLE $m(X, \cdot)$ [UTILITY]

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$$(A \oplus B) \oplus \text{VAL}_i(X, t_i)$$

||

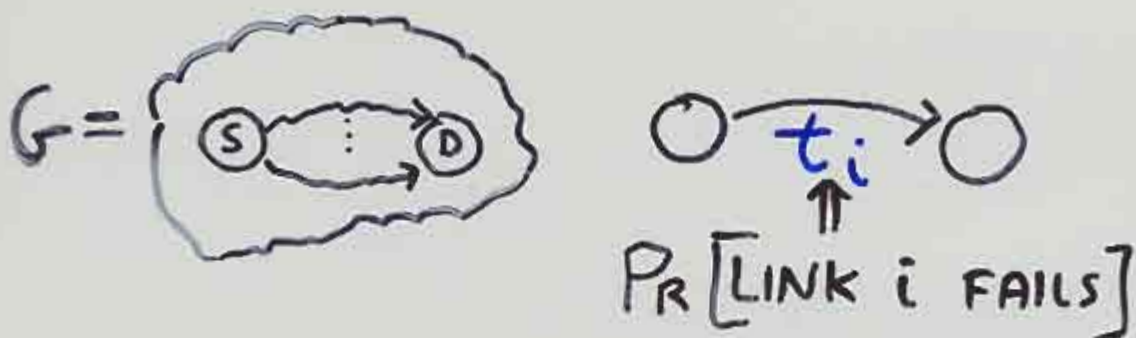
$$A \oplus \text{VAL}_i(X, t_i) \oplus B$$

CONSISTENT: "⊕" ASSOCIATIVE, COMMUTATIVE PROBLEMS

$$\text{SOL}(\mathcal{I}) = \text{SOL}(\pi)$$

↑
PUBLIC PART

MOST RELIABLE PATH



GOAL: FIND MOST RELIABLE
PATH FROM S TO D

MAXIMIZE $m(X, t)$

$$\prod_{i \in X} t_i$$

UTILITY:

EXPECTED PAYMENT $P_i \cdot t_i$

PAYMENTS:

$$P_i = P_R[X \text{ DOES NOT FAIL} \mid i \text{ DOES NOT FAIL}]$$

TASK SCHEDULING

m MACHINES, e JOBS

ASSIGNMENT:

$$X: [e] \longrightarrow [m]$$



$$VAL_i(X, t_i) = - \left(\begin{array}{l} \text{FINISH TIME} \\ \text{OF MACHINE } i \end{array} \right)$$

GOAL: MINIMIZE MAKESPAN

$$\text{MAX}_i \{ -VAL_i(X, t_i) \}$$

CONSISTENT?

YES: "RENT THE MACHINES"

RENTING:

1) PAY α - FINISH TIME

2) NO MORE THAN P_i



MONEY TO AG_i :

$$\text{MIN} \{ \alpha + \text{VAL}_i(x, t_i), P_i \}$$

$$m(x, t) = \text{MIN}_i \{ \text{VAL}_i(x, t_i) \}$$

CONSISTENT! 😊

NOTE: WITHOUT "RENTING" NO
MECHANISM (EVEN FOR $m=2$)

[NISAN-RONEN '99]

RENTING:

- 1) PAY α - ~~FINISH~~ ^{RELEASE} TIME
- 2) NO MORE THAN P_i



MONEY TO AG_i :

$$\text{MIN} \{ \alpha + \text{VAL}_i(x, t_i), P_i \}$$

$$m(x, t) = \text{MIN}_i \{ \text{VAL}_i(x, t_i) \}$$

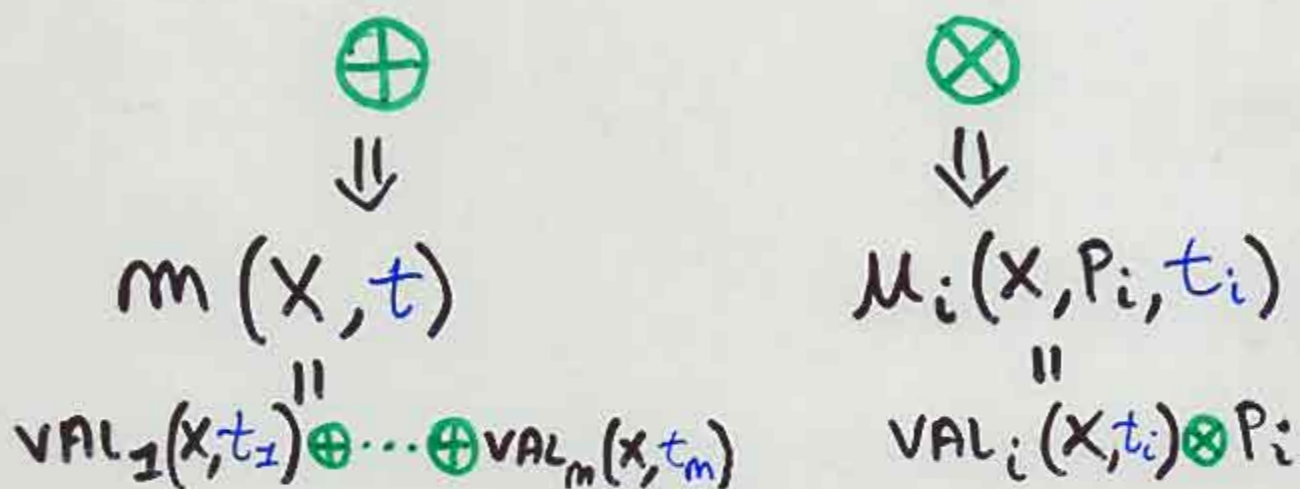
CONSISTENT! 😊

NOTE: WITHOUT "RENTING" NO MECHANISM (EVEN FOR $m=2$)

[NISAN-RONEN'99]

OPEN QUESTIONS

SEMI-CONSISTENT PROBLEMS?



HOW MUCH COMPUTATIONAL ISSUES HELP?

- TASK SCHEDULING
(MAKE THE PROBLEM CONSISTENT)
- SELFISH KNAPSACK
(WHICH INFORMATION ARE "CRITICAL")

INTUITION

OUR GOAL

AG_i's GOAL

$$m(x, t) = \mu_i(x, P_i, t_i)$$

AG_i IS WILLING TO HELP US
(MAXIMIZE $m(x, t)$) ☺

POSSIBLE "=" IF $m(\cdot)$ AND $\mu_i(\cdot)$
HAVE A COMMON "ROOT"

